

# AspectC Language Specification

## Version 0.4 \*

Weigang Gong and Hans-Arno Jacobsen

Middleware Systems Research Group  
Department of Computer Science &  
Department of Electrical and Computer Engineering  
University of Toronto

January 3, 2007

— **Working Technical Draft.**  
For updates and changes, please refer to [www.AspectC.net](http://www.AspectC.net).

---

\*This research is supported by NSERC.

# Contents

<b>1</b>	<b>Overview</b>	<b>4</b>
<b>2</b>	<b>Change Log</b>	<b>4</b>
<b>3</b>	<b>Join Point Model</b>	<b>5</b>
<b>4</b>	<b>Pointcut</b>	<b>5</b>
4.1	Primitive Pointcut . . . . .	5
4.2	Composite Pointcut . . . . .	8
4.3	Named Pointcut . . . . .	8
4.4	cflow() Pointcut . . . . .	10
4.5	Matching Mechanism . . . . .	11
4.5.1	Simple Character Matching . . . . .	12
4.5.2	Wildcard Character Matching . . . . .	12
<b>5</b>	<b>Advice</b>	<b>13</b>
5.1	Single Advice . . . . .	13
5.2	Proceed() . . . . .	14
5.3	“this”: Reflective Information at Join Points . . . . .	15
5.4	Multiple Advice . . . . .	17
5.4.1	before/after advices . . . . .	17
5.4.2	around advices . . . . .	17
<b>6</b>	<b>Static Crosscutting</b>	<b>20</b>
6.1	intype() Pointcut . . . . .	20
6.2	introduce() Advice . . . . .	21
<b>7</b>	<b>Implementation</b>	<b>22</b>
7.1	Aspect Compilation . . . . .	23
7.2	Syntax Analysis . . . . .	23
7.3	Advice Weaving . . . . .	23
<b>8</b>	<b>Grammar</b>	<b>25</b>
8.1	Keywords . . . . .	25

8.2	Grammar Rules . . . . .	25
<b>9</b>	<b>Usage</b>	<b>26</b>
9.1	General Usage . . . . .	26
9.2	Command Line Options . . . . .	27

# 1 Overview

ASPECTC is an implementation of aspect-oriented programming (AOP) for the C programming language. ASPECTC is an extension to C.

This specification introduces the ASPECTC programming model and the new language constructs.

The specification is implemented by the ASPECTC compiler that weaves code written in ASPECTC into ASPECTC-unaware ANSI-C code, and generates C sources implementing the aspect-oriented program. These sources can be compiled by any ANSI-C compliant compiler such as gcc.

The current ASPECTC language design adapts the ideas of aspect-oriented programming laid-out in the original paper by Kiczales *et al.* [2] to the C programming language. The ASPECTC language loosely follows the ASPECTJ programming language design [3] and the partial ASPECTC language design originally suggested by Coady *et al.* [1].

To this end ASPECTC aims at being enabling technology. It is the necessary “evil” and investment in building research infrastructure to eventually lead to further explorations and investigations not possible today, as no stable ASPECTC implementation exists.

Long-term research objectives of the AspectC project include the investigation of

1. concern separation support and aspect-oriented language features tailored to the C language and the imperative style of programming
2. aspect-orientation in the context of software written in C, especially systems software and middleware systems, targeting small-scale, embedded systems (e.g., cell phones, PDAs, chip cards, sensor boards etc.)
3. techniques and tools for the development of highly customizable and easily configurable systems and middleware systems software product lines catering to the extensive world of C-based systems.

The ASPECTC implementation and further information can be found on the project Web pages at [www.AspectC.net](http://www.AspectC.net).

This document is neither a tutorial on ASPECTC, nor a research paper, it simply documents our current ASPECTC implementation and offers the corresponding language specification.

However, this document should suffice to get you started developing ASPECTC programs and sending us bug reports. The ASPECTC distribution contains a lot of test cases illustrating the use of ASPECTC constructs. The project Web pages — [www.AspectC.net](http://www.AspectC.net) — also contain a few examples.

## 2 Change Log

Changes in ASPECTC Version 0.4 relative to ASPECTC Version 0.3:

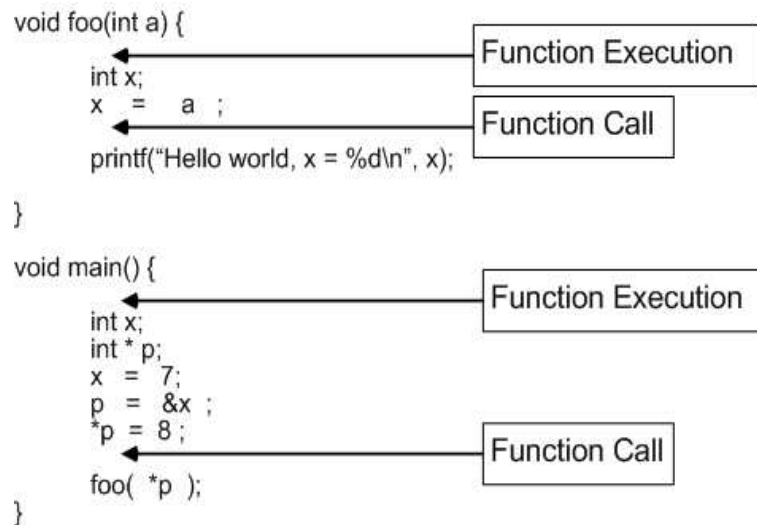
- Section 4: Added `callp()` pointcut description.
- Section 8: Added new keyword “`callp`”.
- Section 9.1: Added a new example to show ASPECTC compilation with pre-processing step.

### 3 Join Point Model

A *join point* is a well-defined point in the execution context of a program. Currently, ASPECTC supports the following join points:

1. call join point: the point when a function is called
2. execution join point: the point when a function is executed

Both join points are illustrated by the program below:



### 4 Pointcut

A *pointcut* is a language extension representing one or more join points. Currently, ASPECTC supports *primitive pointcuts*, *composite pointcuts*, and *named pointcuts*.

#### 4.1 Primitive Pointcut

A *primitive pointcut* represents either one of the two join points defined above. The following pointcuts are defined.

1. `call(function-signature)`

A `call` pointcut picks out the join points of calling the function specified by *function-signature*.

2. `callp(function-signature)`

A `callp` pointcut picks out the join points of calling the function specified by *function-signature* through dereferencing a function pointer.

3. `execution(function-signature)`

An `execution` pointcut picks out the join points of executing the function specified by *function-signature*.

4. `args(a list of types or identifiers)`

An `args` pointcut picks out call or execution join points whose parameters' types match the specified types or the types of the specified identifiers.

5. `infile("file name")`

An `infile` pointcut picks out call or execution join points which appear in the file specified.

6. `infunc(identifier)`

An `infunc` pointcut picks out call or execution join points which appear in the function specified.

7. `result(type or identifier)`

A `result` pointcut picks out call or execution join points whose return type matches the specified type or the type of the specified identifier.

## Semantics:

1. The *function-signature* must be a valid prototype of a function. It represents the function associated with the call or execution join point.
2. For the `callp` pointcut, the function name specified in the signature can not contain the wildcard character.
3. The `callp` pointcut captures function calls by dereferencing the following types of function pointers: global, local and function pointers passed as argument.

For example, "`callp(void foo(int))`" captures the function calls shown below.

```
void (*gp)(int);
struct A {
    void (*sp)(int);
};
```

```

void foo(int a) {
    ...
}
void foo2(void (*ap)(int)) {
    void (*lp)(int);
    (*ap)(); <— capture
    lp = foo;
    (*lp)(); <— capture
    {
        void (*llp)(int);
        llp = foo;
        (*llp)(); <— not capture
    }
}
int main() {
    struct A sa;
    sa.sp = foo;
    (*sa.sp)(); <— not capture
    gp = foo;
    (*gp)(); <— capture
    foo2(foo);
}

```

4. For each parameter in the prototype, only its type can be specified.

For example, “`call(int foo(int ))`” picks out any call to function “`foo`” accepting an `int` parameter and returning an `int`.

5. The identifiers specified in the args or result pointcuts must be declared as a parameter in the advice declaration. The main usage of args and result pointcuts is to expose program context to advice functions.

For example, “`before(int x): args(char, x)`” picks out any call or execution join point whose parameter types are “`char`” and “`int`”, and the value of the second parameter is available for use inside the advice function, as follows:

```

before(int x) : execution ( void foo ( char , int ) ) && args(char , x) {
    printf(“inside before advice, param = %d\n”,x );
}

```

6. There is a special format of args(): “`args(* pointer-variable-name)`”. The meaning is that the parameter type of a join point must match the type after dereferencing

the pointer variable.<sup>1</sup> Using this format, advice functions can change the value of the argument passed into a function, as follows:

```
before(int * x) : execution ( void foo (char c, int n) ) && args(char , *x) {
    *x = (*x) * 2;
    printf("inside before advice, argument value is doubled\n");
}
```

7. The file name specified in the `infile` pointcut must be enclosed by quotes, and it should be the name of a generated file, not the input file.

For example, say the input main file is `t1mc.mc`, if a developer wants to pick out all join points appearing in this file, she must use `infile("t1mc.c")`, because `"t1mc.c"` is the name of the generated file.<sup>2</sup>

8. The identifier specified in the `infunc` pointcut should be the function name where the join point occurs.

## 4.2 Composite Pointcut

A *composite pointcut* defines a pointcut by composing pointcuts with the following operators: `"&&"`, `"||"`, `"!"` or `"()"`. The syntax is as follows:

1. `pointcut0 && pointcut1`: returns join points picked up by both `pointcut0` and `pointcut1`
2. `pointcut0 || pointcut1`: returns join points picked up by either `pointcut0` or `pointcut1`
3. `!pointcut0`: returns join points not picked up by `pointcut0`
4. `(pointcut0)`: returns join points picked up by `pointcut0`

### Semantics:

1. The pointcuts connected by the afore-mentioned operators can be any valid pointcut declaration.

## 4.3 Named Pointcut

To improve usability of pointcuts, developers can attach a name to a pointcut description, and the name can then be used in places where a pointcut is used. The syntax for attaching a pointcut name is as follows:

---

<sup>1</sup> `"before(int *x):args(x)"` is not the same as `"before(int *x): args(*x)"`. The first matches a join point whose parameter type is `"int *"`, but the latter matches a join point whose parameter type is `"int"`.

<sup>2</sup>In future releases we plan to relax this point.

pointcut *pointcut-name* ( *parameter-list* ): *pointcut-description*;

The syntax for using a named pointcut is as follows:

*identifier* ( *identifier-list<sub>opt</sub>* )

For example, the following example shows how to declare a named pointcut and use the name in two different advices:

```
pointcut callFoo() : call(void foo ());
before() : callFoo() && infunc(main) {
    printf("before calling foo in function main\n");
}
before() : callFoo() && infunc(foo2) {
    printf("before calling foo in function foo2\n");
}
```

### Semantics:

1. The *pointcut-name* can be any valid identifier.
2. The *parameter-list* can be empty, indicating there are no exposed arguments associated with the pointcut.
3. The *pointcut-description* can be any valid pointcut.
4. The *identifier* must be the name of a named pointcut.
5. The number of identifiers in the *identifier-list* must be the same as the number of parameters in the *parameter-list* where the named pointcut is declared.
6. The type of each identifier in the *identifier-list* must be the same as that of the corresponding parameter in the *parameter-list*.
7. The name of each identifier in the *identifier-list* should be declared as a parameter of the corresponding advice or named pointcut, such as:

```
pointcut FirstNamedPC(int z) : call(void foo (int)) && args(z);
before(int j) : FirstNamedPC(j) { ... }
pointcut SecondNamedPC (int w) : FirstNamedPC(w) ;
```

8. The developer can also expose the arguments or the return value by using a named pointcut, as follows:

```

pointcut callFoo(int w) : call(void foo (int )) && args(w);
before(int k) : callFoo(k) && infunc(main) {
    printf("before calling foo in function main, value = %d\n", k);
}
before(int p) : callFoo(p) && infunc(foo2) {
    printf("before calling foo in function foo2, value = %d\n", p);
}

```

## 4.4 cflow() Pointcut

ASPECTC provides a `cflow()` pointcut to pick out all join points occurring in the dynamic execution context, i.e., the control flow, of other join points. Its syntax is `cflow(pointcut-definition)`.

For example, “`call(void foo(int)) && cflow(execution(void foo3()))`” only picks out the calls to function `foo` under the control flow of function `foo3`.

Given the following advice:

```

void around() : call (void foo(int)) && cflow(execution(void foo3())) {
    printf("skip call of foo in control flow of foo3\n");
}

```

If the advice is applied to the following C code:

```

void foo(int a) {
    printf("in foo\n\n");
}
void foo2() {
    printf("in foo2\n");
    foo(3);
}
void foo3() {
    foo2();
}
int main() {
    printf("call foo in main\n");
    foo(9);
    printf("———\n");
}

```

```

    printf("call foo2 in main\n");
    foo2();
    printf("-----\n");
    printf("call foo3 in main\n");
    foo3();
}

```

The output is:

```

call foo in main

in foo

-----

call foo2 in main

in foo2

in foo

-----

call foo3 in main

in foo2

skip call of foo in control flow of foo3

```

### Semantics:

1. The pointcut definition inside `cflow()` can be any valid pointcut definition except another `cflow()` pointcut.

## 4.5 Matching Mechanism

ASPECTC provides two mechanisms for matching pointcuts with join points – simple character matching and wildcard character matching.

### 4.5.1 Simple Character Matching

When a plain string is specified in a pointcut's declaration, ASPECTC uses simple case-sensitive string comparison for matching.

For example:

1. `call(int foo(int))` picks out any call to function "foo" accepting an `int` parameter and returning an `int`.
2. `args(int, char)` picks out any call or execution of functions accepting an `int` and a `char` as parameters.
3. `call(int foo(int)) && infunc(foo2)` picks out any call of function "foo" inside function "foo2".

### 4.5.2 Wildcard Character Matching

ASPECTC uses "\$" and "..." as wildcard characters to enhance the matching capability: \$ matches any type identifier or any length of continuous strings, including the empty string; ... matches any length item list, including the empty list.

For example:

1. `call(i$t f$oo(in$))` picks out any call to functions which have a name starting with "f" and ending with "oo", have a return type starting with "i" and ending in "t", and accept one parameter having a type starting with "in".
2. `args(int, ..., char)` picks out any call or execution of functions accepting an `int` and a `char` as the first and last parameters.
3. `call(int foo(int)) && infunc(fo$o2)` picks out any call of function "foo" inside functions whose name starts with "fo" and ends with "o2".

#### Semantics:

1. Developers can use \$\$ to match one \$ inside a target name.
2. ... can only be used when specifying parameter types for a function's prototype.
3. When the name specified in `args()` or `result()` pointcuts has \$, ASPECTC searches an advice parameter having the exact same name. If found, the name is bound with the advice parameter, otherwise, ASPECTC treats the name as a type name.

For example,

- (a) “before(int x\$x): args(char, x\$x)” picks out any call or execution join point whose parameter types are ”char” and ”int”, because “x\$x” matches an advice parameter having type “int”.
- (b) “before(): args(char, x\$x)” picks out any call or execution join point whose first parameter type is ”char” and second parameter type’s name starts with “x” and ends with “x”.

## 5 Advice

### 5.1 Single Advice

An advice represents the code to be executed when a join point is matched by a pointcut defined inside the advice declaration. Currently, ASPECTC supports the following types of advices:

1. before: code is executed before some join points
2. after: code is executed after some join points
3. around: code is executed instead of code at some join points

The general syntax for an advice declaration is:

```
type-specifieropt before|after|around ( parameter-type-listopt ) : pointcuts
{ function-body }
```

#### Semantics:

1. For before/after advice, the *type-specifier* should not be specified. The ASPECTC compiler uses “void” as the return type of the function generated from the advice.
2. For around advice, the *type-specifier* must be specified, and it becomes the return type of the function generated from the advice. Furthermore, the type-specifier must be the same as the return type of the matched functions.
3. The ASPECTC compiler generate a unique function name for each advice.
4. If the *parameter-type-list* is specified, it becomes the parameter list of the generated function.
5. The information specified by the *pointcuts* is used to match join points.
6. For each parameter name in the *parameter-type-list*, the name must be used inside one pointcut among the *pointcuts*, like args() or result().

For example:

```

before() : execution ( void foo (int ) ) {
    printf("before execution foo\n");
}

```

The “before” advice indicates that a message is printed out before the execution of function “foo”.

```

int around() : call ( int foo (int ) ) {
    printf("around call foo\n");
    return 100;
}

```

This “around” advice indicates that a message is printed out, 100 is returned, and the calling of function “foo” is skipped.

```

after(int k) : execution ( void foo (int ) ) && args(k) {
    printf("before execution foo, argument = %d\n",k);
}

```

This “before” advice takes a parameter which exposes the argument value of function foo to the advice function.

## 5.2 Proceed()

Around advice can be used to skip code at an existing join point. However, sometimes developers still want to access the original function call/execution inside the advice. This can be achieved by using `proceed()` inside the around advice. The `proceed()` call takes the original value of the arguments<sup>3</sup> and calls/executes the original function.

For example:

```

int around() : call ( int foo (int) ) {
    printf("around call foo\n");
    printf("value of foo = %d\n", proceed());
    return 0;
}

```

This shows that function `foo()` is accessed inside an around advice, and its return value is used by the advice.

---

<sup>3</sup>Note, the developer can use the `args()` construct to change the value of the original argument. If `proceed()` is called afterward, the new value is used to call/execute the original function.

### 5.3 “this”: Reflective Information at Join Points

Within the advice code, ASPECTC provides a special pointer variable, `this`, to access reflective information about the current join point.<sup>4</sup>

The following string fields can be accessed by `this`:

1. `funcName`: the function name of the join point.
2. `kind`: the join point kind, either “call” or “execution”.

For example, the following advice prints out the name of every called/executed function in the control flow of `main`:

```
before(): cflow(execution(int main())) {
    printf("aspect: kind = %s, function = %s\n", this->kind, this->funcName);
}
```

When the advice is applied to the following C code:

```
void foo(int a) {
    printf("in foo\n\n");
}
void foo2() {
    printf("in foo2\n");
    foo(3);
}
void foo3() {
    foo2();
}
int main() {
    foo3();
}
```

The output is:

---

<sup>4</sup>`this` is similar to the `thisJoinPoint` variable in ASPECTJ.

aspect: kind = execution, function = main

aspect: kind = call, function = foo3

aspect: kind = execution, function = foo3

aspect: kind = call, function = foo2

aspect: kind = execution, function = foo2

in foo2

aspect: kind = call, function = foo

aspect: kind = execution, function = foo

in foo

## 5.4 Multiple Advice

When a join point is matched by pointcuts from multiple advices, the various types of advices are handled differently.

### 5.4.1 before/after advices

Advices are executed sequentially according to the matching sequence.

For example:

```
/* advice 1 */
before() : execution ( void foo (int ) ) {
    printf("before advice 1");
}

/* advice 2 */
before() : execution ( void foo (int )) {
    printf("before advice 2");
}
```

Since the execution join point of function `foo` is matched by both advice 1 & 2 and both are before advices, the two advices are executed in sequence. That is advice 1 is executed before advice 2.

### 5.4.2 around advices

- `no proceed()`: the first matched advice is executed, and the rest are skipped.
- `has proceed()`: the `proceed()` call inside the around advice invokes the next matched around advice if there is one; otherwise, the `proceed()` call invokes the original function.

For example:

```
/* advice 3 */
void around() : execution ( void foo (int) ) {
    printf("around advice 3");
}
```

```

/* advice 4 */
void around() : execution ( void foo (int)) {
    printf("around advice 4");
}

```

The execution join point of function `foo` is matched by both around advices 3 & 4. Since there is no `proceed()` inside the advices, the first matched advice, advice 3, is executed, and advice 4 is skipped.<sup>5</sup>

The situation changes, if a `proceed()` call is present in the advice. For example:

```

/* advice 5 */
void around() : execution ( void foo (int) ) {
    printf("around advice 5");
    proceed() ;
}

```

```

/* advice 6 */
void around() : execution ( void foo (int)) {
    printf("around advice 6");
    proceed() ;
}

```

Since there is a `proceed()` call used inside the advices, the execution sequence is: advice 5 → advice 6 → `foo`.<sup>6</sup>

By using multiple around advice and `proceed()`, the developer can impose different advices for join points. This can achieve effects similar to multiple if-statements, like:

```

/* advice 7 */
void around(int x) : execution ( void foo (int)) && args(x) {
    if(x < 3) {
        printf("around advice 7");
        return;
    }else {

```

---

<sup>5</sup>Even if there is a `proceed()` call inside advice 4, it is not executed, since the execution join point is already surrounded by advice 3 without `proceed` (i.e., defined as “around” advice of advice 3 that does not let the call `proceed` to either further advice or the surrounded code).

<sup>6</sup>If there is no `proceed()` in advice 6, the original function `foo()` will not be executed.

```

        proceed();
    }
}

/* advice 8 */
void around(int x) : execution ( void foo (int)) && args(x) {
    if(x < 9) {
        printf("around advice 8");
        return;
    }else {
        proceed();
    }
}

/* advice 9 */
void around(int x) : execution ( void foo (int)) && args(x) {
    if(x < 20) {
        printf("around advice 9");
        return;
    }else {
        proceed();
    }
}

```

The effects of applying advice 7, 8, & 9 is same as: whenever calling a function “foo” with parameter “x”,

```

if(x < 3) {
    printf("around advice 7");
}else if(x < 9) {
    printf("around advice 8");
}else if(x < 20) {
    printf("around advice 9");
}else {
    foo(x);
}

```

## 6 Static Crosscutting

In addition to expressing dynamic crosscutting represented by call and execution join points, ASPECTC also provides mechanism to support static crosscutting, such as the addition of members to structs and unions.

### 6.1 `intype()` Pointcut

An `intype` pointcut picks out the struct or union type whose name matches the type name specified or which has been typedefed by a name matching the one specified. Its syntax is “`intype(identifier)`”.

For example, if the types are declared as follows:

```
struct X {           <— first struct
    int a;
};
typedef struct X MYX1;
typedef MYX1 MYX2;
typedef struct {     <— second struct
    int b;
} MYX3;
typedef MYX3 MYX4;
```

then the `intype()` pointcut has the following effects:

1. “`intype(struct X)`”, “`intype(MYX1)`”, and “`intype(MYX2)`” match the first struct.
2. “`intype(MYX3)`” and “`intype(MYX4)`” match the second struct.
3. “`intype(MYX$)`” matches both structs.

#### Semantics:

1. The *identifier* must be a struct or union type name, or a name assigned by a typedef for a struct or a union.
2. The wildcard character `$` can be used inside the *identifier*.
3. The `intype()` pointcut must only be used within an `introduce()` advice.

## 6.2 introduce() Advice

An `introduce()` advice adds new data members to the struct or union type picked out by the `intype()` pointcut. Its syntax is:

```
introduce (): pointcuts { function-body }
```

### Semantics:

1. The *pointcuts* must contain only `intype()` pointcuts, or contain composite or named pointcuts built from `intype()` pointcuts.
2. The *function-body* must contain valid struct or union member declarations.
3. The data member name declared inside the *function-body* must not collide with the existing member names in the matched type.
4. ASPECTC simply copies the *function-body* and adds them to the end of the matched struct or union declaration.
5. If multiple `introduce()` advices are applied to the same type, the data members from each advice are added according to the matching sequence.

For example, for the types described above and the following advice declarations:

```
introduce() : intype(struct X) {                <— advice 1
    double b;
}
introduce() : intype(MYX1) {                    <— advice 2
    double c;
}
introduce() : intype(MYX3) {                    <— advice 3
    double c;
}
introduce(): intype(MYX3) || intype(MYX1) {     <— advice 4
    char * p;
}
```

After the advices are applied, the types become:

```
struct X {
    int a;
```

```

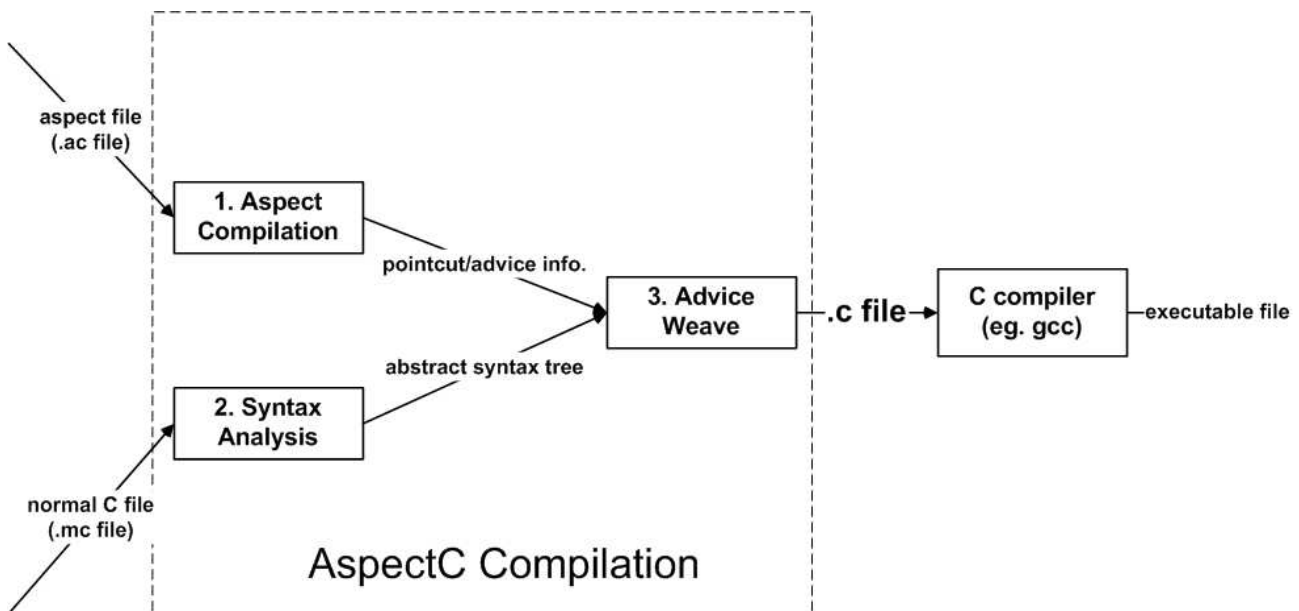
double b;           <— from advice 1
double c;           <— from advice 2
char * p;           <— from advice 4
};
...
typedef struct {
    int b;
    double c;       <— from advice 3
    char * p;       <— from advice 4
} MYX3;
...

```

## 7 Implementation

ASPECTC is implemented as a source-to-source translator. The inputs are ASPECTC files and C source files<sup>7</sup>. The aspect files contain pointcut, advice, or normal C code. The outputs are normal C files with advice code inserted at the point specified by pointcuts. The output files can then be compiled by a C compiler.

There are 3 phases in the ASPECTC compilation process: aspect compilation, syntax analysis and advice weaving. The compilation process is described by the following figure.



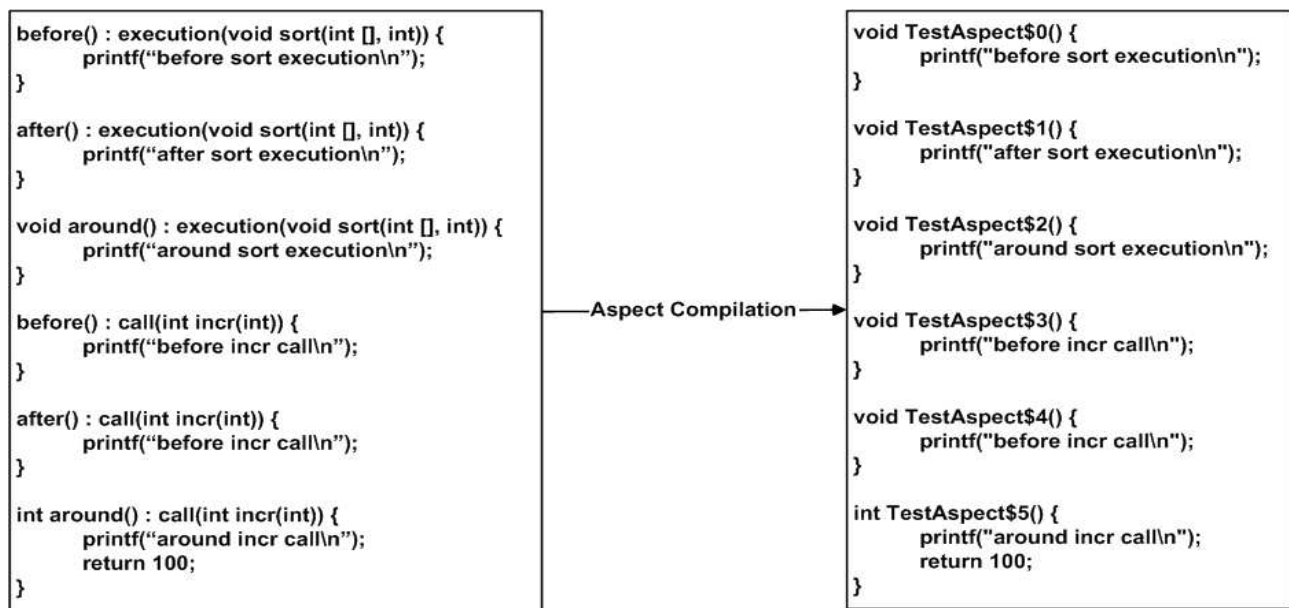
<sup>7</sup>Both kinds of input files need to be pre-processed by a C pre-processor or by gcc using the "-E" option.

## 7.1 Aspect Compilation

In the aspect compilation phase, each advice is compiled to a unique C function. The advice parameters are compiled to parameters of the new C function. In the advice weaving phase, these parameters are bound to function arguments. Since before advice and after advice have no return type, the ASPECTC compiler uses “void” as return type for the corresponding functions. For around advice, the ASPECTC compiler uses the return type specified in the advice declaration as the return type of the function.

Another task in this phase is to collect information related to pointcut and advice, which is used in the advice weaving phase.

The following figure illustrates the kind of C functions generated from the advice in the aspect file.



## 7.2 Syntax Analysis

The main purpose of this phase is to collect information to facilitate join point matching by generating an abstract syntax tree (AST) for the C sources.

## 7.3 Advice Weaving

The last phase is to insert calls to advice functions in appropriate locations in the C sources. The following figure illustrates how calls are inserted into a C file.

```

before() :
    execution(void sort(int [], int)) {
        printf("before sort execution\n");
    }
after() :
    execution(void sort(int [], int)) {
        printf("after sort execution\n");
    }
void around() :
    execution(void sort(int [], int)) {
        printf("around sort execution\n");
    }
before() : call(int incr(int)) {
    printf("before incr call\n");
}
after() : call(int incr(int)) {
    printf("before incr call\n");
}
int around() : call(int incr(int)) {
    printf("around incr call\n");
    return 100;
}

```

```

void sort(int x[], int n) {
    printf("here is sort\n");
}
int incr(int x) {
    x = x + 1;
    return x;
}
int main() {
    int x[5] = {3,5,2,1,4};
    int a;
    sort(x,5);
    a = 8;
    a = incr(a);
    return 0;
}

```

Advice Weave →

```

void sort(int x[], int n) {
    { TestAspect$0(); }
    { TestAspect$2(); }
    { TestAspect$1(); }
    return ;
}
int incr(int x) {int $rtValue;
    (x=(x+1));
    $rtValue = x;
    return $rtValue;
}
int incr$SelectionSort$0 (int x ) ;
int main() {int $rtValue;
    int x[5] = {3,5,2,1,4};
    int a;
    sort(x, 5);
    (a=8);
    (a=incr$SelectionSort$0(a));
    $rtValue = 0;
    return $rtValue;
}
void TestAspect$3 () ;
void TestAspect$4 () ;
int TestAspect$5 () ;
int incr$SelectionSort$0 (int x ) {
    int $rtValue;
    { TestAspect$3(); }
    { $rtValue = TestAspect$5();}
    { TestAspect$4(); }
    return $rtValue;
}

```

## 8 Grammar

In order for ASPECTC to support aspect-oriented language extensions, the following keywords and grammar rules are added to the C language grammar.

### 8.1 Keywords

New keywords are : “args”, “after”, “around”, “before”, “call”, “callp”, “execution”, “infile”, “infunc”, “introduce”, “intype”, “pointcut”, “proceed”, “cflow”, “result”.

### 8.2 Grammar Rules

*function-definition:*

*declaration-specifiers<sub>opt</sub> declarator : pointcuts compound-statement*

*declaration:*

**pointcut** *declarator : pointcuts ;*

*pointcuts:*

*or-pointcuts*

*pointcuts && or-pointcuts*

*or-pointcuts:*

*unary-pointcut*

*or-pointcuts || unary-pointcut*

*unary-pointcut:*

*base-pointcut*

**!** *base-pointcut*

*base-pointcut:*

**args** ( *type-or-id-list* )

**call** ( *func-jointpoint* )

**execution** ( *func-jointpoint* )

**infile** ( *string-literal* )

**infunc** ( *identifier* )

**intype** ( *type-name* )

**result** ( *type-or-id* )  
*identifier* ( *identifier-list<sub>opt</sub>* )  
**cflow** ( *pointcuts* )

*func-jointpoint*:  
*declaration-specifiers declarator*

*type-or-id-list*:  
*type-or-id*  
*type-or-id-list type-or-id*

*type-or-id*:  
*type-name*  
*identifier*

*direct-declarator*:  
**before**  
**after**  
**around**  
**introduce**

## 9 Usage

### 9.1 General Usage

The ASPECTC compiler takes C source files with and without ASPECTC syntax as input. The files with ASPECTC syntax must have the suffix ".ac" and those without AspectC syntax must have suffix ".mc". Furthermore, both types of files should be pre-processed by a C preprocessor before passing through the ASPECTC compiler.<sup>8</sup>

The ASPECTC compiler outputs ANSI-C compliant C source files to be processed by a C compiler. If any input file has an unknown suffix, the ASPECTC compiler emits an error message and stops compilation.

Example 1:

Suppose there are neither #include nor macro directives used in the a.ac or the b.mc files:

```
>acc a.ac b.mc
```

---

<sup>8</sup>In the future we plan to relax the suffix restriction for C files.

The ASPECTC compiler generates `a.c` and `b.c` C-source files for processing by a C compiler, like for example `gcc`.

```
>gcc a.c b.c
```

```
>./a.out
```

Example 2:

Suppose there are `#include` or macro directives used in the `a.ac` or the `b.mc` files. This requires that the source files have to be pre-processed before weaving and compilation.

1. since `gcc` does not recognize the `.ac` or the `.mc` suffix, the suffix must to be changed to `.c`.

```
>cp a.ac a_ac.c
```

```
>cp b.mc b_mc.c
```

2. pre-process the files by a pre-processor, and save the output in files with the by the ASPECTC compiler required suffixes

```
>gcc -E a_ac.c > a_ac.ac
```

```
>gcc -E b_mc.c > b_mc.mc
```

3. perform the ASPECTC compilation

```
>acc a_ac.ac b_mc.mc
```

4. the ASPECTC compiler generates `a_ac.c` and `b_mc.c` C-source files for processing by a C compiler, like for example `gcc`.

```
>gcc a_ac.c b_mc.c
```

```
>./a.out
```

Makefile examples capturing the above steps are part of the ASPECTC distribution.

## 9.2 Command Line Options

The following command line options are supported by the ASPECTC compiler:

1. -t

-thread-safe

The code generated to support the `cflow()` pointcut is thread-safe.<sup>9</sup>

2. -v

-version

The compiler's Version number is printed.

## References

- [1] Yvonne Coady, Gregor Kiczales, Mike Feeley, and Greg Smolyn. Using `aspectc` to improve the modularity of path-specific customization in operating system code. In *ESEC/FSE*, 2001.
- [2] G. Kiczales, J. Lamping, A. Mendhekar, C. Maeda, C. Videira Lopes, J.-M. Loingtier, and J. Irwin. Aspect-oriented programming. In *ECOOP*, 1997.
- [3] AspectJ Team. AspectJ project web site. <http://www.eclipse.org/aspectj/>.

---

<sup>9</sup>The ASPECTC compiler uses a GCC specific feature, `-thread-local storage`, to ensure thread-safety for the generated code. The `__thread` keyword is used. However, since this is not a standard feature and not supported by all C compilers, ASPECTC turns the option off by default. The default code generated for `cflow` pointcuts is not thread-safe. For more information about the thread-local storage feature, visit <http://gcc.gnu.org/>.